

1 OPERATING SYSTEMS LABORATORY XIII

UNIX FILE MANAGEMENT

- In UNIX, each file is represented by an inode
- The inode
 - maps individual byte addresses relative to the beginning of the file to logical block numbers for a particular disk
 - holds permission information for the file, whether the file is a regular file, a directory, a special file, and the current file size
- An example inode

```
-----  
|      block #      |  
-----  
|      block #      |  
-----  
|      block #      |  
-----  
|      block #      |  
-----  
|      block #      |  
-----  
|      block #      |  
-----  
|      block #      |  
-----  
|      block #      |  
-----  
| permissions, ownership, |  
| current file size, file |  
| type                  |  
-----
```

- blocks are a single, fixed size,
- table index corresponds to logical position in the file

- For example,

	block offset in file
----- block 34	0
----- block 722	1
----- block 1072	2
----- block 6	3
----- block 377	4
----- block 771	5
----- block 7	6
----- block 83	7
----- block 212	8
----- block 433	9
----- block 812	single
----- block 96	double
----- block 531	triple
----- permissions, ownership, etc	

defines a file. If blocks are 512 bytes long, the 1033rd byte in the file is found by first calculating the logical offset (in blocks) from the beginning of the file as:

$$1033 / 512 = 2$$

At table entry 2 (block offset 2) we find logical disk block number 1072. The byte offset within disk block 1072 is:

$$1033 \% 512 = 9$$

So given the inode shown above, the 1033rd byte of the file is byte 9 of block 1072 on the disk from whence this inode is allocated.

-- latter table slots (single, double, triple in the figure) refer to single-, double-, and triple-levels of indirection.

-- the number of direct, single, double, and triple slots in each inode are implementation specific. However, given these numbers, the block size, and the size of a block number, it is possible to calculate how large the largest file that can be represented is.

For example, assume

block number = 4 bytes
block size = 512 bytes
direct blocks = 10
single indirect blocks = 1
double indirect blocks = 1
triple indirect blocks = 1

implies

(10) directly accessible blocks +
(512 / 4) single indirectly addressable blocks +
(512 / 4)² double indirectly addressable blocks +
(512 / 4)³ triple indirectly addressable blocks =

2113674 addressable blocks * 512 bytes/block =
1,082,201,088 addressable bytes within a single file

- \$ ls -li (execute this command inside your home directory)

- Compile and run the following program

- Analyze the output

/* structuremembers.c

Print structure members of files

st_mode the type and mode of the file

st_ino

st_dev

st_rdev

st_nlink

st_uid

```
st_gid
st_size
st_atime
st_mtime
st_ctime
*/
#include <stdio.h>
#include <sys/types.h>
#include <sys/stat.h>

int main()
{
    struct stat status;
    int i;
    for (i=1 ; i < argc ; i++)
    {
        if(stat(argv[i], &status))
            fprintf(stderr, "Cannot stat %s \n", argv[i]);
        else
            printf("%15s %4.4o\n", argv[i], status.st_mode & 07777);
            //printf("%15s %14d\n", argv[i], status.st_ino);

    }

    return 0;
}
```